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A Survey of Distributed Dataset Synchronization in Named Data Networking

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***Abstract*—Distributed synchronization of a shared dataset (sync for short) provides a powerful abstraction for connection-agnostic multiparty communication in NDN. In recent years, several sync protocols have been proposed, each featuring different design choices in data naming, namespace representation, and state propagation mechanisms, which lead to different design tradeoffs. In this report, we survey these protocols and, through detailed analysis and side-by-side comparisons, highlight their commonalities and fundamental differences. We also articulate the remaining issues that must be addressed to make the sync protocols available to more applications, shedding the light on future work directions in this important area of NDN architecture research.**

I. Introduction

Named Data Networking (NDN) [1], [2] is a proposed new Internet architecture that shifts the communication model from host-centric, as in today's TCP/IP networks, to data-centric. At the network layer, NDN provides a simple yet powerful communication primitive that allows a data consumer to send an Interest packet with a name or name prefix and retrieve a Data packet which is named under that prefix and can be verified. While the Interest-Data exchange primitive has significantly narrowed the semantic gap between the application layer and the network layer in today's TCP/IP network architecture, it is cumbersome to use directly to build distributed applications that often involve some form of data or state sharing and synchronization among multiple parties. For example, file sharing, collaborative editing, and group messaging all collect and distribute state and data among groups of participants. However, with the TCP/IP architecture, whenever communication involves more than two parties, the applications have to either establish multiple TCP connections between the peers or rely on (at least logically) centralized infrastructure to support multiparty communication.

The data-centric nature of the NDN architecture provides a foundation for distributed dataset synchronization (*sync* for short) as an important layer of abstraction for reliable multiparty communication on top of the Interest-Data exchange primitives. Distributed applications and services that use sync can easily publish data in a shared dataset maintained by the sync protocol without worrying about how others would discover the newly available data, and can easily consume up-to-date information without worrying about how and from where to get it. This use of NDN sync to provide reliable data-centric communication differs from data retrieval via a TCP connection in three fundamental ways. First, it naturally supports data retrieval among multiple parties, while TCP supports data exchange between two parties only. Second, it does not require all communicating parties to be interconnected at the same time as TCP does. Third, it does not care from where the data is returned since the security is attached to the data instead of its container or communication channel.

NDN can achieve distributed dataset synchronization by synchronizing the *namespace* of the shared dataset among a group of distributed nodes (called *sync nodes*), thanks to the unique and secured binding between name and content. To share new data, a producing (side of) applications injects its names into the dataset. After learning the new names, the consumer (sides of) application decides whether to fetch the new data according to its own needs and available resources. One may view sync as playing a *transport* layer role in the NDN architecture, bridging the gap between the functionality required by the distributed applications and the one-Interest-one-Data datagram retrieval semantics offered by NDN network-layer primitives.

In this report, we present a survey of the distributed data synchronization protocols that have been developed for the NDN architecture in recent years. They include CCNx 0.8 Sync [3], iSync [4], CCNx 1.0 Sync [5], ChronoSync [6], RoundSync [7], and PSync (or PartialSync) [8]. We identify three most important questions in a sync protocol design, and examine each of the above protocols by how it answers these design questions, to extract common design patterns as well as differences in the approaches. We conclude this survey with a discussion on a set of identified remaining issues.

II. Existing Sync Protocols Over NDN

In this section, we examine the set of existing sync protocols that have been developed for the NDN architecture. Our goal is to extract common design patterns for NDN sync protocols and identify different design choices and tradeoffs made in different protocols. Our analysis focuses on the following key design aspects:

*a) **Data naming:*** Thanks to the unique binding between names and immutable data object in NDN, a shared dataset can be uniquely identified by the namespace containing the hierarchical names of all data packets in the dataset. Therefore the dataset synchronization problem in NDN is conveniently reduced to the synchronization of the corresponding namespace. The sync protocol may directly synchronize the application data names, or leverage the sequential naming convention to simplify the dataset namespace (and encapsulate the application data names if necessary).

*b) **Namespace representation**:* The data structure that represents the state of the shared dataset namespace is often referred to as the *sync state*. Every sync node keeps a local copy of the sync state and uses the sync protocol to keep up with the changes generated by other nodes in the sync group. This requires the sync state to encode the namespace without loss of information and allow sync nodes to detect and reconcile the differences in the shared namespace between distinct states.

*c) **State sync mechanism**:* Each node participating in a sync group may publish new data in the shared dataset at any time. The sync protocol should ensure the other nodes in the group can receive the new data and reach agreement on the state of the dataset. The state synchronization mechanism therefore needs to enable the nodes to (1) learn about the updates as soon as possible and (2) detect and reconcile inconsistency in the sync state caused by other factors such as packet loss and/or network partition.

To support dataset synchronization inside a group, the sync protocol also requires a group communication namespace for the sync nodes to publish and exchange protocol messages. To achieve group communication, the protocol may rely on the underlying network to provide multicast capability (which is the case in all existing sync protocols), or explore other group rendezvous mechanisms such as structured communication (e.g., Distributed Hash Table [9]) and epidemic dissemination [10]. Note that the design of the group communication mechanism is outside the scope of the sync protocol (and therefore not the focus of our analysis), but may have a significant impact on the protocol design choices. Section III-A discusses the group rendezvous issue and different sync communication models in more detail.

We summarize the commonalities and differences among the six protocols in Table I. In the rest of this section, we describe each sync protocol by focusing on the three design aspects mentioned above. At the end of the section we give a brief summary that highlights common design patterns and provide preliminary comparison on the efficiency issues including synchronization delay and protocol message size. (An extended evaluation of existing sync protocols will be reported in future revisions.)

A. CCNx 0.8 Sync

The CCNx 0.8 Sync protocol [3] is the earliest synchronization solution proposed for the NDN/CCN architecture as a service module of the *ccnr* repo daemon. CCNx 0.8 Sync allows a set of repos to synchronize a shared data collection that contains data with arbitrary application names. The set of data names under a common *collection* prefix is organized into a tree structure called the *sync tree* (see Fig. 1), where a node in that tree may store a single data name (i.e., a *leaf*) or a list of (leaf and non-leaf) nodes. The structure of the sync tree is determined by the order in which the data names are added to the collection, which is independent from the canonical ordering of the data names.

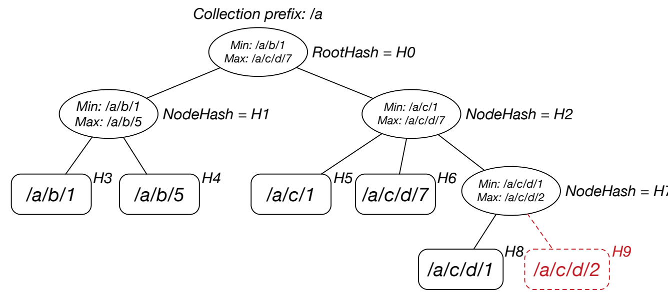

Fig. 1: Example of a sync tree in CCNx 0.8 Sync

Each node in the sync tree is associated with a hash value: the value of the leaf node is simply the hash of the name stored in that node; the value of the non-leaf node is recursively computed as the arithmetic sum of the hashes of all its children. In other words, the hash value of a node is the sum of the hashes of all data names contained in the sub-tree under that node. For example, in Fig. 1, $H_3 = Hash(/a/b/1)$, $H_2 = H_5 + H_6 + H_7$, and $H_0 = H_1 + H_2$. The root hash (H_0 in Fig. 1) then provides a summary of the entire namespace (i.e., sum of all data name hashes). Note that the sum of hashes is not a cryptographically strong summary: in certain cases two sync trees may store different sets of names but happen to have the same root hash.

Any producer connected to a repo can publish new data into the data collection at any time. The sync module in the repo daemon (called *sync agent*) keeps track of the insertions of new data and updates the sync tree accordingly, adjusting the hash values along the path from the new leaf node to the root. For example, in Fig. 1 the insertion of a new data "`/a/c/d/2`" (marked as the red dashed square at the bottom right) will cause the sync agent to update the node hashes H_7 and H_2, eventually propagating the change up to the root hash H_0.

The sync agent periodically advertises the latest root hash by sending a *RootAdvice* Interest to all the other repos that store the same data collection. The RootAdvice Interest name starts with a multicast prefix for the sync tree, which is shared by all repos and different from the collection prefix, followed by the current root hash of the sync tree. When a sync agent receives a remote root hash that is different from its own, it replies to this RootAdvice with its own root hash. The sync agent who receives a RootAdvice reply will send a *NodeFetch* Interest, which is also named under the multicast prefix of the sync tree, to the replying repo to retrieve the list of hashes for all the children under the root node of the remote sync tree. The NodeFetch process is recursively applied to all the nodes in the sync tree, skipping those with the same hash value between local and remote, until all nodes with different hash values have been visited. Once it learns the names of the new data from the leaf nodes, the sync agent can fetch those data from the remote repo via normal Interest-Data exchange and insert that data to its local copy of the data collection. An example of the synchronization process in CCNx 0.8 Sync (triggered by the update to the sync tree shown in Fig. 1) is illustrated in Fig. 2. Note that while we show the sync protocol messages only between two repos for clarity, the RootAdvise

TABLE I: Comparison of existing sync protocols in NDN

	CCNx 0.8 Sync	iSync	CCNx1.0 Sync	ChronoSync	RoundSync	PSync
Synchronized Namespace	Arbitrary names	Arbitrary names	Arbitrary names	Node prefix + seq#	Node prefix + seq#	Stream prefix + seq#
Sync state representation	Hash tree	IBF of hashes of names	Manifest storing names or digests of data	List of {prefix : seq#}	List of {prefix : seq#} + round log	IBF of hashes of names with highest seq#
State change detection	Data replying to *RootAdvice* Interest with local root hash	Interest carrying digest of IBF	Interest carrying hash of manifest	Data replying to *Sync Interest* with updates	*Sync Interest* carrying digest of current round	Data replying to *Sync Interest* with new IBF
State update retrieval	*NodeFetch* Interest retrieving child node hashes	Interest retrieving IBF content	Interest retrieving manifest	Data replying to *Sync Interest* with updates	Data replying to *Data Interest* with updates in current round	Data replying to *Sync Interest* with new IBF

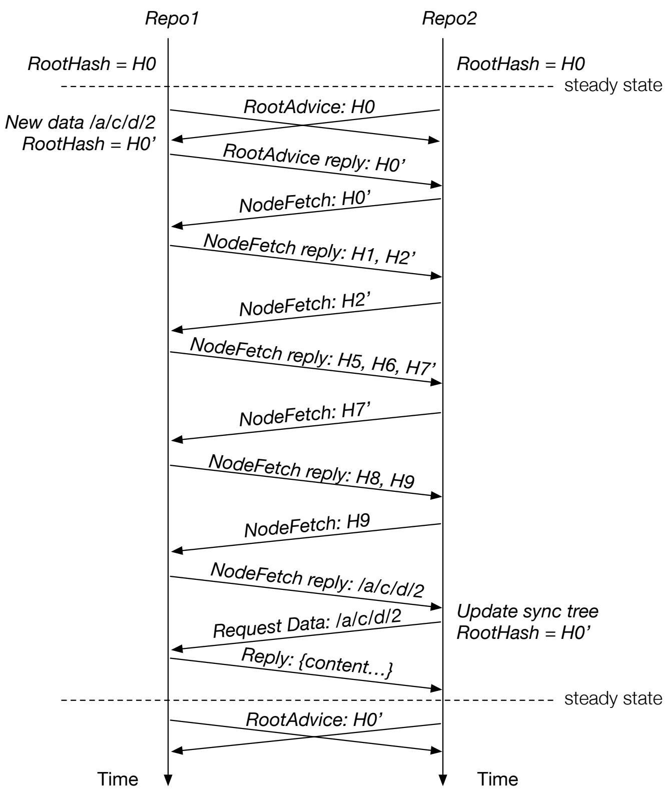

Fig. 2: Synchronization in CCNx 0.8 sync

and NodeFetch Interests actually carry multicast prefix and will be received by all repos storing the same data collection.

One problem in the update propagation mechanism in CCNx 0.8 Sync is that when multiple repos publish new data simultaneously, there will be more than one reply to a RootAdvice Interest and only one of them will be returned to the Interest issuer. In such case, the sync agent who sends the initial RootAdvice Interest need to issue additional Interests to fetch other replies. The proposed solution is to attach *exclude filters* to the Interest to list the root hashes of the remote sync trees that have already been received. This ensures that each unique remote sync tree is examined only once for missing data.

A side-effect of the CCNx 0.8 Sync algorithm, which compares the local and remote sync trees and updates the local state to be the union of the two, is that the repo cannot remove any data once it is added to the data collection. This is because the algorithm cannot distinguish the case where a repo intentionally removed a piece of received data from the case where the repo has never received the data before. As a result, the data collection maintained by CCNx 0.8 Sync must be monotonically growing, which creates usability issues with the applications who generate a large mount of data and need to perform garbage collection periodically to reclaim the storage. For example, when the NDNVideo application [11] was deployed on top of CCNx repo to publish live video streams, the system administrator had to cleanup the data and restart all repo instances every day at midnight in order to avoid overwhelming the storage of the repo server.

B. *iSync*

iSync [4] is a direct optimization on top of the CCNx 0.8 Sync design. Like in CCNx 0.8 Sync, it supports the synchronization of shared data with application names. To represent the sync state more efficiently, iSync uses Invertible Bloom Filter (IBF) [12] to store all the names from the shared dataset in compressed form. Since the IBF can only store fixed-length items, the data names must be first mapped to fixed-length IDs (generated from the hash of the names) before they are added to the IBF. A bi-directional mapping table is maintained by every sync node so that it can recover the original NDN names from the IDs.

Different from CCNx 0.8 Sync,[1] iSync uses "digest broadcast" Interests (equivalent to the RootAdvise Interest in CCNx) to advertise its current state to other nodes periodically, rather than a solicitation for different sync states. Since the encoded size of the IBF is typically very big, the advertisement Interest carries only the digest of the current IBF from the sending node. When a node receives a digest different from its own, it sends another Interest to request the corresponding IBF content. After it receives the IBF from a remote node, the

[1]The original iSync paper [4] describes the CCNx 0.8 Sync protocol differently compared to the official specification [3] released in the CCNx source code package.

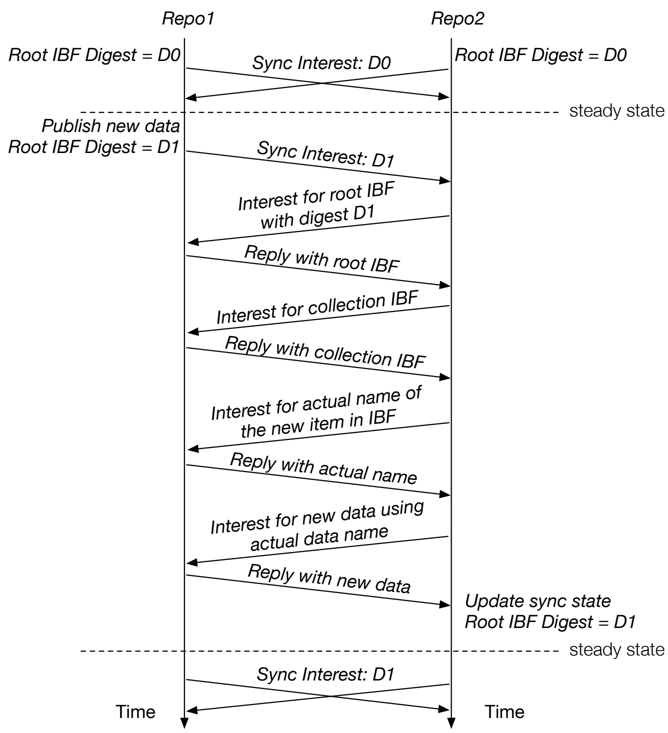

Fig. 3: Synchronization process in iSync

node subtracts its own IBF from the remote IBF and extracts individual IDs from the resulting "diff" IBF. Once the sync node extracts all new IDs, it issues Interests to request the original NDN names corresponding to those IDs and then fetches the new data using the original names. An example of the synchronization process in iSync is shown in Fig. 3.

Note that the iSync node does not expect any reply to the initial advertisement Interest it sends. It therefore gets around the issue with multiple replies generated for the same broadcast or multicast Interest. Note that this causes imbalance in the Interest-Data packet flow in the network. Assuming the number of advertisement Interests is negligible compared to the total network traffic, the impact can be a tolerable amount of waste of PIT entries that eventually expire and get removed.

A major limitation in the IBF data structure is that it can losslessly encode up to a certain number of items, beyond which some of the stored items cannot be extracted. Unfortunately, it is not uncommon that during the synchronization process the set difference between the namespace of some sync nodes may contain too many IDs that cannot be encoded in the IBF in a lossless fashion. iSync provides several ways to control the size of the set difference at multiple levels in the protocol design. First, the shared dataset is divided into multiple collections that host data for different applications; each collection maintains its own IBF independently from others. Second, iSync protocol enforces each node to periodically advertise its local sync state and resolve the difference, which bounds the delay of the data propagation and the size of the set difference between any two nodes. Third, iSync creates multiple *local IBFs* to record the small-step changes during each sync period; if the advertised IBF (called *global IBF*) contains too many changes, the sync node can fetch the local IBFs instead and perform more fine-grained difference reconciliation.

C. CCNx 1.0 Sync

The design proposal of CCNx 1.0 Sync [5] abandons the CCNx 0.8 Sync design and adopts a simple manifest-based solution.[2] The manifest packets are named under a routable *data collection prefix* announced by every sync node, followed by the hash of the manifest and segment numbers. The manifest contains the SHA256 hashes or the exact names of all data objects in the shared data collection. When the SHA256 hashes are used, the names of the data objects are constructed by appending the hash value to the same data collection prefix in the manifest name. The application-layer data (with real application names) may be encapsulated in those data objects.

Each sync node uses Interest packets to advertise the hash of its local catalog manifest when it generate new data. The advertisement Interests are also named under the data collection prefix and forwarded to all sync nodes announcing that prefix. They have short lifetime and do not retrieve any data. To increase the possibility that all nodes can receive the advertisement, the node repeats the advertisement Interest once or twice within a few seconds after the first advertisement is sent. Once a node receives a different hash, it should also advertise its own hash under the control of some gossip protocol (with random backoff and duplicate suppression). It then sends out Interests to retrieve the corresponding (possibly segmented) manifest packets, compares the names listed in the manifest with its local namespace, and then retrieves the missing data over the network. This approach is similar to iSync but without the benefit of efficient encoding and differentiation provided by the IBF data structure.

D. ChronoSync

ChronoSync [6] attempts to improve efficiency of dataset synchronization by utilizing naming conventions. In particular, each ChronoSync node publishes data that contain application-layer messages under its own unique name prefix. This prefix also serves as an identifier for the node in the sync group and is aligned with the topological prefix of the access network for each node. The name of the data is constructed by concatenating the node prefix with a sequence number that starts from zero and gets incremented by one for each new data published by the sync node. Although in theory the sequence number could wrap around if represented by fixed-size integers, in practice it is usually not a big issue: for a sync node publishing data at the rate of 1000 packets per second, it takes more than half a billion years for a 64-bit sequence number to wrap around; even with 32-bit integer representation it still takes about 50 days for the sequence number to wrap around, which provides enough time for the sync group to garbage-collect the previous data.

The sync node maintains a 2-level "flat" sync tree, as is shown in Fig. 4, with each leaf containing the data prefix

[2]The authors are not aware of any real implementation of this new design.

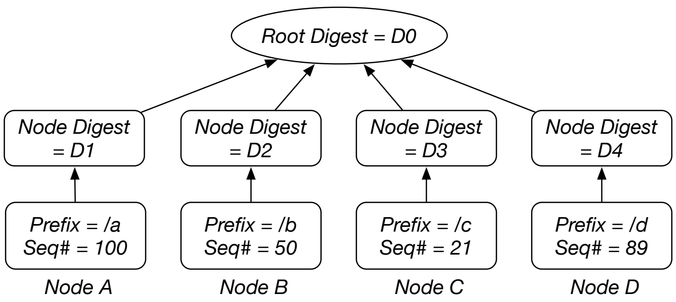

Fig. 4: Example of a sync tree in ChronoSync

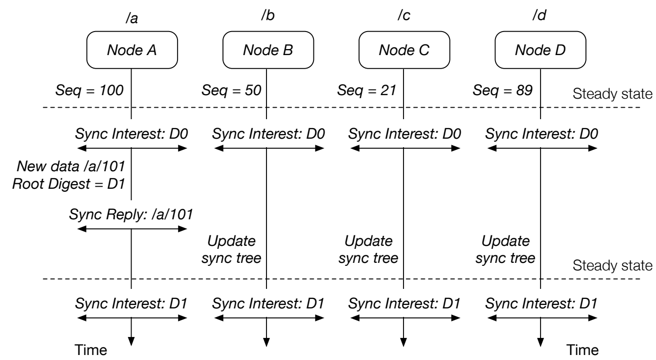

Fig. 5: Synchronization process in ChronoSync

and the latest sequence number of each producer in the sync group. Each leaf is associated with the digest calculated over node's prefix and the latest sequence number. The root of the tree maintains the digest of concatenation of leaf digests canonically ordered by the corresponding prefix names. Since the naming convention is to publish data with continuously increasing sequence numbers (starting from zero), this sync tree is essentially a condensed representation of the namespace containing all the data ever published in the group, and the root digest is a short summary of the dataset.

ChronoSync nodes maintain *long-lived Sync Interests* in the network by transmitting a new Sync Interest immediately when the previous one expires or gets satisfied. The long-lived Interest stays in the pending Interest table of the forwarders in the network so that any reply to the Sync Interest can be returned to every node in the group as soon as it is generated. The Sync Interest name starts with the multicast sync group prefix and carries the current root digest of the sender's local sync tree. The Sync Interest serves two important purposes: first, it advertises the sender's digest in the group so that other nodes can detect inconsistency in the sync state; second, it solicits the next state changes generated on top of the state identified by the digest carried in the Sync Interest.

In the steady state, all nodes generate identical state digests and send out the same Sync Interest that is aggregated by the NDN forwarders. When some node publishes new data and increments its sequence number, instead of replying to the long-lived Sync Interest with its new root digest as in CCNx 0.8 Sync, the node replies with the name of its newly published data (i.e., the node prefix and the sequence number).[3] This *Sync Reply* is efficiently delivered to all the other nodes in the group, following the multicast tree built by the pending Sync Interest. After they receive the reply, the nodes update their local sync tree, recompute the root digest, and then send out Sync Interests carrying the new digest. An example of the synchronization process in ChronoSync is shown in Fig. 5.

To allow efficient state reconciliation, each ChronoSync node maintains a limited log of historical digests and the corresponding dataset states. If some node is lagging behind in the synchronization process and sends out a Sync Interest with a digest that has been observed by other nodes, these sync nodes can respond with all the data published in the group since that digest is announced. Note that when multiple sync nodes reply to the Sync Interest carrying a previous digest (potentially with different sets of updates if they are not synchronized), at most one of those relies will be received by the sender of that Interest. Nevertheless, the reply helps speed up the synchronization process of the Interest sender who is trying to catch up with the rest of the group.

There are several cases where a node may receive Sync Interests with unrecognized digests. In the first case, a node may receive a Sync Interest with an updated digest before receiving the Sync Reply that triggered the update. To handle that situation, ChronoSync injects a random delay to process the Sync Interest with unknown digest at a later time, expecting to receive the corresponding Sync Reply while waiting.

In the second case, multiple Sync Replies can be generated in response to the same Sync Interest, if multiple nodes publish new data at the same time. However, because of NDN's flow balance property, nodes will receive no more than one reply to the Sync Interest. As a result, nodes may receive different data items, compute multiple different state digests, and start announcing them in the sync group.

The third and a more complicated case arises if the network is partitioned for a long period of time and then reconnected. The sync nodes in different partitions have cumulated multiple updates to the sync tree, leading to a sequence of digests that are unrecognizable to the nodes in other partitions.

ChronoSync can handle simple cases when the nodes diverge by at most one Sync Reply by resending the previous Sync Interest with exclude filters that contain the implicit digests of the received Sync Replies. However, if multiple changes have been applied to the sync state at some node, the mechanism using exclude filters will not be able to retrieve the diverging sync replies generated by every node (see II-E for detail). In such cases, ChronoSync falls back to a *recovery* mechanism: when a node observes an unknown digest, it will send out a special *Recovery Interest* containing the unknown digest; the nodes who recognize that digest will reply with the complete information about its sync tree, rather than the specific changes that lead to that digest; when the requesting node gets the reply, it will merge the received sync tree into its local sync tree by taking the higher sequence number from both trees for each sync node.

To support more complex naming conventions with richer

[3]If multiple data packets are generated, the Sync Reply carries only the largest sequence number of all new data.

embedded semantics (e.g., for trust management), application can use ChronoSync together with "one level of indirection." In other words, application-defined data names (and the data itself if the size is small) can be encapsulated in the data packets managed by the sync layer. This can be realized either directly or with the help of sync-managed *actions* that describe objects added or removed from the application data collection, such as updated files a distributed file system or messages in a chat room.

E. RoundSync

The recognition of the ChronoSync problem in scenarios with many simultaneous data generation led to development of the RoundSync [7] protocol. Specifically, one of the causes of the problem is the overloaded function of Sync Interests: (1) to detect different states among the sync nodes and (2) to retrieve the updates from other nodes. As a result, the Sync Replies carrying the updates to the shared dataset will be named after the previous Sync Interest name which contains the digest of the corresponding sync state. If a node generates Sync Replies on top of a diverged state (e.g., in the scenario with partitioned sync group), nodes with different states will not be able to derive the name for those Sync Replies and therefore cannot send Interests to retrieve them. Merging the diverged sync states will only create new set of sync states, potentially contributing in further divergence of the states. To re-synchronize in this case, ChronoSync must rely on a recovery mechanism to receive the entire sync state.

To address this problem, RoundSync divides the synchronization process into *rounds*, updates semantics of the Sync Interest, and introduces the new type of Interest packet called *Data Interest*. The RoundSync's Sync Interest, augmented with the round number information, serves only as a notification mechanism (similar to iSync) to inform other sync nodes about the state in the round. When the divergence is detected, the nodes can request the data in the round using the predictably named Data Interests, i.e., names of Data Interests do not include state digest but only the round number. Therefore, published data within a specific round can be retrieved even if the states are not fully synchronized. The replies to the Data Interest have the same functionality as the Sync Reply in the original ChronoSync design, i.e., they carry the node's prefix and sequence number of the newly published data. In addition, RoundSync mandates that a sync node can publish at most one data packet in each round and must move to a new round when it receives new data published by others in the current round. This helps reduce the chances of state divergence caused by simultaneous data production.

For example in Fig. 6, a sync node may start publishing data at round 11 even though it is still trying to synchronize with other nodes at round 10 or earlier. If multiple nodes publish data in the same round simultaneously, they will detect the inconsistency through Sync Interest and then send Data Interests with exclude filters to retrieve those Data Interest replies. Since there will be at most one reply from each node in a single round, the exclude filter mechanism will allow the nodes to eventually retrieve all updates.

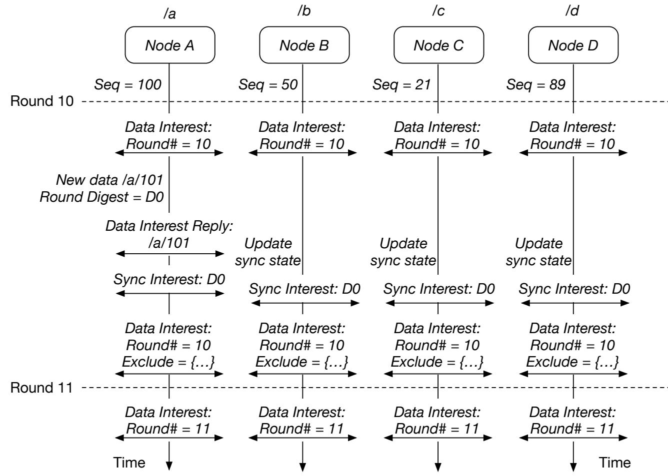

Fig. 6: Synchronization process in RoundSync

RoundSync maintains digest for each round in a *rounds log* table. To allow nodes who missed the Sync Interests in earlier rounds to detect and recover the missing data, RoundSync also computes *cumulative digests* that covers the entire dataset as observed in a round and is piggybacked in the Data Interest replies of future rounds. Upon receiving a different cumulative digest for some round that is long before the node's current round, the sync node sends out a Recovery Interest to fetch the full sync state and the current round number S from the node who generated that cumulative digest, instead of retrieving missing data round-by-round (which may take a long time). After receiving the reply, the node merges the received dataset with its own, discards the rounds log entries for the rounds before S and resumes normal RoundSync operation for the rounds after S.

F. PSync

PSync [8] (a.k.a. PartialSync) was originally designed for the consumers to synchronize a subset of a large data collection with a single producer. The data packets published by the producer are organized into *data streams* which are identified by the unique stream prefixes. Like in iSync, PSync also employs IBF to represent the namespace by storing the hash of the names (called *KeyID*) in the fixed-length slot of the IBF. However, PSync also adopts the naming convention in ChronoSync and RoundSync that data packets from the same stream are ordered by the continuous sequence numbers. Therefore the IBF only needs to store the latest data name from each stream. This further reduces the amount of information stored by the IBF and allows the applications to choose a smaller IBF size that can be transmitted more efficiently over the network.

To support the synchronization of a subset of the producer's data (a.k.a., *partial sync*), PSync introduces the subscription list to encode the prefixes of the data streams that the consumer is interested in.[4] The subscription list is a Bloom Filter (BF)

[4]PSync allows the consumers to specify their subscription only at the granularity of data streams.

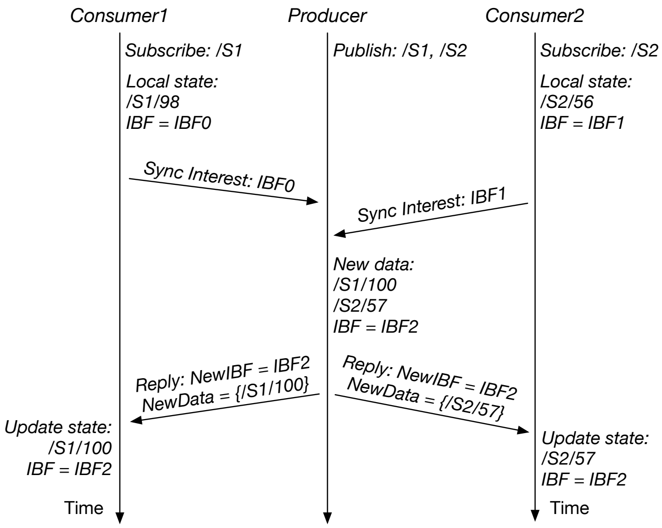

Fig. 7: Synchronization process in PSync

that stores the hashes of those stream prefixes. The size of the Bloom Filter is determined by the total number of streams a consumer may subscribe to and the false positive rate the consumer is willing to accept. Special cases like empty and full subscription may be encoded more efficiently with special markers.

During the sync process, the consumer keeps a local copy of the producer's IBF which indicates the data it has received so far. To sync up with the producer and retrieve new data, the consumer maintains *long-lived Sync Interest* whose name contains the local IBF copy and the consumer's subscription list. When the producer publishes new data, it first subtracts the IBF in the pending Sync Interest from its new IBF, and extracts the KeyIDs of the new data packets that have not been received by the consumer yet. Then the producer checks whether the stream prefixes of those new data packets are included in the consumer's subscription list (subject to certain false positive rate). Finally the producer generates a Sync Reply containing the original names of the new data packets in the subscribed streams and also its latest IBF. Upon receiving the Sync Reply, the consumer updates its local IBF copy with the received IBF, and sends out Interests to fetch the new data. An example of the synchronization process in PSync is shown in Fig. 7.

An important feature in the PSync design is that each consumer maintains its own data consumption and subscription status. The producer, on the other hand, does not maintain per-consumer state, which significantly reduces the amount of data stored by the producer. If multiple producers are serving the same set of data streams, the consumers may send Sync Interests via *anycast* to get replies from any producer that is available online, assuming that these producers have run sync protocols among themselves to sync up their dataset. However, this *stateless* producer design introduces two additional costs: first, the Sync Interest and Sync Reply need to carry the IBF and the subscription list (BF) which will bloat the size of the Interest name up to hundreds of bytes; second, the producer needs to generate Sync Replies in real-time for each Sync Interest since it does not remember the previous consumption status of each consumer and cannot pre-generate the next Sync Reply.

Although it was initially designed for producer-consumer synchronization, PSync can support multi-producer distributed dataset synchronization (like other sync protocols discussed in this paper) where each sync node is both producer and consumer at the same time. This is achieved by having every sync node subscribe to all data streams published by every other node. However, in this "full synchronization" mode each node only needs to maintain a single IBF which represents the state of the whole dataset, rather than keeping a separate IBF for each node in the group. In addition, the Sync Interests need to be forwarded via *multicast* to the entire group so that any node who has produced new data can respond with a reply that carries the updates.

G. Summary

From the previous discussion, we can see that a few common design patterns have arisen in the key design aspects among the existing NDN sync protocols. ChronoSync, RoundSync, and PSync have adopted the sequential data naming conventions, i.e., naming the data packets using sequence numbers under a common name prefix for each producer or data stream, to simplify the representation of the shared dataset namespace. Having continuous and monotonically increasing sequence numbers in the data name allows the cumulative data collection generated by the same producer or in the same data stream to be summarized by the highest sequence number. This reduces the amount of information that needs to be encoded in the sync state and also simplifies the protocol design since the sync protocol needs to only focus on synchronizing the latest sequence numbers rather than the whole namespace.

The existing sync protocols have used a variety of data structures to represent the sync state. All of those data structures provide lossless encoding of the data names (or the hashes of the names) in the shared dataset. CCNx 0.8 Sync and CCNx 1.0 Sync enumerate the dataset namespace in the hierarchical sync tree and the manifest, respectively. To reconcile the set difference, the sync nodes simply compare the content in the sync tree or the manifest and then retrieve the missing data from the remote nodes. ChronoSync and RoundSync also enumerate the dataset namespace by listing the latest sequence numbers from all data producers in the sync state. State reconciliation is achieved by comparing the sequence number of each producer between the local and remote sync states and taking the maximum as the latest sequence number. iSync and PSync use IBF to compress the dataset namespace and perform set reconciliation using IBF subtractions. However, due to the limited IBF capacity, both iSync and PSync have to provide means for controlling the size of the set difference between different nodes.

The existing sync protocols typically use one of the two communication models for propagating the information about the new data published in the sync group. The first model is to use multicast Interests to advertise the summary of the sync state changes (e.g., digest of the updated sync state), which serves as a notification to prompt other nodes in the sync group to retrieve detailed information about the changes. The second model is to have the sync nodes send "long-lived" Interests to each other (typically using multicast) to pre-establish the return path for the data packet that carries the information about the sync state changes. The "long-lived" Interests essentially become a "one-packet" subscription to the sync state updates generated in the future.

It is often difficult to compare the efficiency of different sync protocols because it usually depends on the application scenarios and the implementation choices. Table II compares the existing NDN sync protocols on a few important performance metrics. One metrics is the data dissemination delay, i.e., the number of round-trips necessary for propagating new data to other nodes. In CCNx 0.8 Sync and iSync, the synchronization process is triggered periodically based on an internal sync timer. Once the process starts, the number of round-trips required to retrieve all updates from a remote node in CCNx 0.8 Sync depends on the depth of the sync tree, while in iSync the process usually finishes within 3.5 RTT, unless the number of changes exceed the capacity of the global IBF in which case the nodes need to retrieve additional "local IBFs". CCNx 1.0 Sync triggers the synchronization process when there is new data published in the dataset, and the data dissemination delay depends on how the nodes retrieve the segmented manifest (e.g., sequentially or pipelined).

ChronoSync and RoundSync achieve optimal synchronization delay when there is no simultaneous data publishing. If multiple nodes generate Sync Replies at the same time, the protocols need additional round-trips to retrieve all Sync Replies using Interests with exclude filters. Therefore the worst-case RTT will be proportional to the number of simultaneous updates in the group, which is bounded by the number of data publishing nodes in the group. PSync achieves the data dissemination delay of 1.5 RTT because the Sync Interests carry specific information about the state of the consumer, which allows the producer to reply with specific changes without spending additional round-trip to request more information. Note that both ChronoSync and PSync require maintaining long-lived Sync Interests in the network so that the replies can be propagated to other nodes as soon as possible. This leads to the overhead of keeping long-lived soft state in the forwarders' PIT (with one PIT entry per sync group).

Another performance metrics is the packet size of the sync protocol messages, which reflects the network bandwidth requirement of the sync communication. Here we mainly focus on the encoding size of the sync state (or state updates) carried in the Interest and/or Data packets. For iSync and PSync, the size of the IBF that summarizes the dataset namespace depends on the size of the hash function output and the data publishing rate of the applications. In a typical implementation that uses 64-bit hash functions and 32-bit counter values, the size of each slot in the IBF is 20 bytes. For an IBF with capacity of 20 items (i.e., allowing at most 20 items to be extracted successfully), the encoded size of the IBF [12] is around 1.5 * 20 * 20 = 600 bytes. Once the size of the IBF is chosen, all Interest and Data packets carrying the IBF will have the same size even if the number of updates is lower than the maximum capacity. Note that PSync usually requires a smaller IBF than iSync because the sequential data naming simplifies the namespace and effectively bounds the number of changes by the number of data prefixes. This enables PSync to carry the IBF directly in the Sync Interest. In iSync, the number of changes within an Interest period depends on the data publishing rate and is unbounded, therefore requiring a larger IBF to accomodate bursty data publishing events.

In contrast, ChronoSync and RoundSync require only the updates to be propagated in reply to the Sync Interests and Data Interests, respectively. Those update data packets contain the prefix and the latest sequence number from each producer who has published new data. Assuming the average size of the data name (i.e., prefix + sequence number) is 40 bytes, the maximum content size of the update is 40 * N, where N is the number of producers in the sync group. In practice, not all producers will be publishing at the same time and the size of the update packets is typically smaller than in the IBF-based approaches. The Sync Interests and Data Interests in ChronoSync and RoundSync usually carry the current state digest only. However, when simultaneous data publishing happens, the nodes need to send additional Interests with exclude filters that enumerate the implicit digests of all the previously received replies. This may cause the size of the Interest packets to grow linearly with the number of simultaneous replies.

In CCNx 0.8 Sync, the size of the NodeFetch reply packets is proportional to the number of children under the requested node in the sync tree; also, the protocol requires multiple NodeFetch packets to resolve all the differences. In CCNx 1.0 Sync, the size of the manifest is proportional to the number of the included data names, representing either a complete shared dataset or serving as (hierarchical) links to manifests (or manifests of manifests) of sub-datasets.

III. Open Issues

In this section we discuss open research issues in distributed data synchronization in NDN that have not been addressed by the existing sync protocols. By inspecting the range of design choices in those open areas, we hope to shed light on the directions for future work.

A. *Group Rendezvous*

The group communication model in the sync protocols has created challenges for the routing scalability in the NDN network, because the sync protocol typically relies on Interest multicast to deliver the state change notification to every node in the group. Supporting Interest multicast via routing would require the per-application multicast prefixes be announced by all networks where the sync nodes reside, which is usually not feasible for large networks hosting many sync groups. One solution currently under our investigation is to utilize a *multicast overlay* that contains a number of dedicated rendezvous points in the network. Those rendezvous points are responsible for collecting and delivering the Sync Interests via the overlay to every sync node. Another solution to the scalable multicast

TABLE II: Comparison of existing NDN sync protocols on common performance metrics

	CCNx 0.8 Sync	iSync	CCNx 1.0 Sync	ChronoSync	RoundSync	PSync
Data dissemination delay	Interest period + tree walk	Interest period + 3.5 RTT (+ RTT to retrieve local IBFs)	Depending on how to retrieve manifest segments	Min is 0.5 RTT; can be long with simultaneous publishing	Min is 1.5 RTT; can be long with simultaneous publishing	1.5 RTT (assuming a single producer)
Interest overhead	Periodic	Periodic	One per update	Long-lived Interest	Two per update	Long-lived Interest
Factors affecting Interest size	Node hash	IBF digest	Manifest digest	State digest (+ exclude filter)	Round digest (+ exclude filter)	IBF + subscription list
Factors affecting Data content size	Number of children under the requested node	IBF size (depending on the number of new data)	Size of the entire dataset	Number of names with new seq#	Number of names with new seq# in a round	IBF size + number of names with new seq#

problem is to establish some communication topology in the sync group, e.g., using *Distributed Hash Table* (DHT) [9], so that the Sync Interests can be propagated via that topology in a peer-to-peer fashion. The third solution is to leverage *viral propagation* (also called *epidemic dissemination*) [10], where each node disseminate new data via NDN sync to its one-hop neighbors (e.g., nodes within the coverage of the wireless signal) and the neighbors further propagate the data until all nodes in the network have received it. The viral propagation model is suitable for disruptive and ad hoc environments where it is infeasible to establish stable communication structure or deploy infrastructural support for multicast.

B. Data management

A sync node usually stores two types of information locally: the data generated by the applications, and the internal state (or metadata) created by the protocol itself. There are several practical issues related to the management of the dataset and protocol state that may have significant impact on the operations of the applications running on top of sync.

1) Data sharing: To improve the data availability, every sync node should be able to serve the entire dataset, including the data published by other nodes, to the consumers. For example, in a group chat application like ChronoChat [13], a user who join the chat room late may want to retrieve some of the earlier chat messages published by other users, even if those users have long left the chat room. In that scenario it is beneficial if the current users in the chat room could store and serve the data published by the past users.

The challenge here is that the application data generated by different peers can be named under the unicast prefixes of different nodes.[5] To fetch the shared data from non-authoritative peers, a sync node can utilize the *Forwarding Hint* [14] mechanism. When a node wants to retrieve some historical data published by a node that is no longer reachable, it may issue Interests with the forwarding hint field carrying one or more unicast network prefixes of the current group members. The NDN forwarders will use the forwarding hint to direct the Interest toward one of the existing nodes in the group.

2) Data archiving: Long-running applications often accumulate a large amount of data. When a new sync node joins the application, it may take a long time to bootstrap the node by fetching the data objects in the shared dataset piece-by-piece. To improve the efficiency of application data transfer, it may be needed to consolidate the data into a single archive file. For example, if the application publishes "actions" through the sync protocol to modify its data (e.g., editing files in a shared folder), the archive file may contain only the latest version of the application data that reflects all the changes made by the users. This mechanism shares a similar spirit with the compaction process in Log-Structured Merge (LSM) databases [15] that coalesces multiple disjoint DB tables into a larger sorted table in order to speed up the lookup operations.

The archive file could be published under the application namespace and segmented into multiple large data packets that can fit into the Maximum Transmission Unit (MTU) of the underlying network. If necessary, the archive data may be stored permanently in dedicated repos without incurring additional storage cost for the sync nodes. Consumers who are interested in the historical data may send pipelined Interests (with the Forwarding Hint as described earlier) to fetch the entire archive file from the sync nodes or the application repos.

3) Data sharding: As the application continues to generate new data, the size of the whole dataset may exceed the storage size of each individual node. If the application requires permanent storage of all data ever published by the users, the shared dataset needs to be *sharded* across the sync nodes for storage scalability. One way to achieve that is to build a data sharding service on top of the sync protocol using consistent hashing [16] or Distributed Hash Table (DHT) techniques to divide the data namespace among the sync nodes. When a sync node receives notification of a new data packet, it consults the data sharding service to decide whether it is responsible for storing that data. If the data falls into its local shard, the node will subsequently retrieve the data and store a local copy; otherwise, it simply updates the local sync state without fetching the data. Note that the sync nodes are still able to sync up with each other even if each of them maintains only a subset of the shared data because the sync state is solely

[5] Having all sync nodes publish data in a shared application namespace (identified by a common multicast application prefix) would require support for Interest anycast in the network.

based on the namespace of the dataset.[6]

C. Group membership management

Having a consistent view of the group membership among the nodes in the distributed system is a prerequisite for a lot of useful functions such as system snapshot, garbage collection, access control, and strong consistency guarantee. For example, to generate a group-wide snapshot of the published data, the sync protocol needs to collect the latest data publishing state from every node in the group; to safely garbage-collect some historical data from the shared dataset, the sync protocol needs to make sure that all nodes in the group have received that data.

Existing sync protocols do not explicitly manage the membership information for the sync group. CCNx 0.8 Sync and iSync do not have the concept of "sync group" at all: the sync state (sync tree and IBF, respectively) does not reflect the identities of the repos that are maintaining the shared dataset. ChronoSync maintains the list of current participants in the sync group via the sync tree, but does not require the sync nodes to have a consistent view of the group while synchronizing with each other. If a sync node receives a new data object produced by a node that is not in its current sync tree, the receiving node will extend its sync tree by adding the producer node into it. Consequently, existing NDN applications running on top of the sync protocols (e.g., ChronoChat) currently have to implement custom group membership management solutions at the application layer.

D. Consistency and data ordering

Consistency in distributed systems has been extensively studied for decades. Strong consistency models such as *linearizability* [17] enforce a global total ordering of events observed by every node in the system. Weaker consistency models relax on the ordering requirements in different ways. In particular, a system with *eventual consistency* is allowed to diverge and expose inconsistent states during the system execution, as long as it eventually resolves the inconsistency.

We must keep in mind two important factors in all consistency discussions: a) the definition of consistency varies among different applications, and b) there exists a tradeoff between consistency and availability. The primary goal of NDN sync is to facilitate multi-party data-centric communication in a distributed system, on top of which the applications can implement different consistency models as application-layer services to meet their semantic requirements. Therefore all existing NDN sync protocols support only the weak consistency model that favors availability over consistency: they allow sync nodes to publish new data at any time and propagate the data to other nodes asynchronously. In the absence of new data generation and permanent network failure, all sync nodes will eventually receive all data packets published by others. Applications with strong consistency requirements can always establish a consistent ordering on the data generated in the sync group using various algorithms and protocols developed in the distributed systems area.

[6]Note that this requires the implementation of the sync protocol to decouple namespace synchronization from data fetching.

IV. Conclusion

This paper presents an overview of the distributed dataset synchronization problem in NDN and a survey on the existing sync protocols. By summarizing and comparing their protocol design, we articulate the different design choices made in the existing sync protocols together with their advantages and limitations. We also discusses open issues that have not been fully addressed in the previous works. By writing this survey paper, we hope that new sync protocols developed in the future can benefit from the past experience and address the open problems with innovative ideas.

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