

Data Structures

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Linear Data Structures

Java's List interface

- ◆ `Get(idx)`
- ◆ `Set(idx, value)`
- ◆ `Add(idx, value)`
- ◆ `Remove(idx)`
- ◆ `listIterator(pos)`

(My)ArrayList Implementation

- ◆ See Figures 3.15 and 3.16 from Weiss text
- ◆ Maintains the following data
 - list of items in an array called `theItems[]`
 - Array capacity (length of above array)
 - Current size called `theSize`
- ◆ Allows the following operations
 - Change in capacity (capacity doubled if array fills up)
 - No change upon removal
 - Implementation of `get(idx)` and `set(idx,x)`
 - Implementation of `size()`, `isEmpty()`, `clear()`
 - Implementation of Iterator interface
 - Index called `current`
 - `next()`, `hasNext()`, `remove()`

add and remove in ArrayList

- ◆ Both involve moving items
- ◆ Operation `add(idx,x)` involves moving all items from position `idx` onward to move in order to make space for `x`
- ◆ Operation `remove(idx)` involves moving all items from position `idx+1` to close the gap created by the removal
- ◆ Study carefully how `ArrayIterator` is implemented

Time Complexity of List Operations

	ArrayList	LinkedList
get(idx)	O(1)	O(N)
set(idx, x)	O(1)	O(N)
add(idx,x)	O(N)	O(N)
remove(idx)	O(N)	O(N)
addBefore(p,x)		O(1)
Remove(p)		O(1)

Stack ADT

```
public interface Stack { // Stack class ADT
    // Reinitialize the stack.
    public void clear();

    // Push "it" onto the top of the stack
    public push(Object it);

    // Remove and return the element at the top of the stack
    public Object pop();

    // Return a copy of the top element
    public Object peek();

    // Return the number of elements in the stack
    public int length();
}
```



Stack Class using Arrays

```
class AStack implements Stack {  
    private Object stackArray[]; // Array holding stack  
    private static final int defaultSize = 10;  
    private int maxSize; // Maximum size of stack  
    private int top; // First free position at top  
  
    // Implementation of all the operations  
}
```


ArrayStack Implementation

```
public boolean push(Object it) {  
    if (top >= maxSize) return false;  
    stackArray[top++] = it;  
    return true;  
}
```

```
public Object pop() {  
    if (top == 0) return null;  
    return stackArray[--top];  
}
```

```
public Object peek() {  
    if (top == 0) return null;  
    return stackArray[top-1];  
}
```

Stack Class using Linked Lists

```
class LinkedStack implements Stack {  
    private Node top;           // Pointer to first element  
    private int size;          // Number of elements  
  
    // Implementation of all the operations  
}
```

LinkedStack Implementation

```
public boolean push(Object it) {  
    top = new Node(it, top);  
    size++;  
    return true;  
}
```

```
public Object pop() {  
    if (top == null) return null;  
    Object it = top.data;  
    top = top.next;  
    size--;  
    return it;  
}
```

```
public Object peek() {  
    if (top == null) return null;  
    return top.data;  
}
```

	ArrayStack	LinkedStack
push(x)	O(1)	O(1)
pop()	O(1)	O(1)
peek()	O(1)	O(1)

Queue ADT

```
public interface Queue { // Queue class ADT
    // Reinitialize queue
    public void clear();

    // Put element on rear
    public boolean enqueue(Object it);

    // Remove and return element from front
    public Object dequeue();

    // Return front element
    public Object frontValue();

    // Return queue size
    public int length();
}
```



Queues using Arrays

```
class AQueue implements Queue {  
    private Object queueArray[]; // Array holding elements  
    private static final int defaultSize = 10;  
    private int maxSize;        // Maximum size of queue  
    private int front;          // Index of front element  
    private int rear;           // Index of rear element  
  
    // Implementation of all the operations  
}
```

Why
front and
rear?

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ArrayQueue Implementation

```
public boolean enqueue(Object it) {  
    if (((rear+2) % maxSize) == front) return false; // Full  
    rear = (rear+1) % maxSize; // Circular increment  
    queueArray[rear] = it;  
    return true;  
}
```

```
public Object dequeue() {  
    if(length() == 0) return null;  
    Object it = queueArray[front];  
    front = (front+1) % maxSize; // Circular increment  
    return it;  
}
```

```
public int length() { return ((rear+maxSize) - front + 1) % maxSize; }
```

Need to distinguish between full and empty.
Hence array is declared to be of size $N+1$
while using only N spots.
Thus, full is checked by:
 $(rear+2) \% maxSize == front$
Empty is checked by: $length() == 0$

Queue using Linked Lists

```
class LQueue implements Queue {  
    private Link front; // Pointer to front queue node  
    private Link rear; // Pointer to rear queue node  
    private int size; // Number of elements in queue  
  
    // Implementations of all operations  
}
```

LinkedList Implementation

```
public boolean enqueue(Object it) {  
    rear.next = new Link(it, null);  
    rear = rear.next;  
    size++;  
    return True;  
}
```

```
public Object dequeue() {  
    if (size == 0) return null;  
    Object it = front.next.data; // retrieve the value  
    front.next = front.next.next; // Advance front  
    if (front.next() == null) rear = front; // Last element  
    size--;  
    return it; // Return element  
}
```

```
public Object frontValue() {  
    if (size == 0) return null;  
    return front.next.data;  
}
```

	ArrayQueue	LinkedList
enqueue(x)	O(1)	O(1)
dequeue()	O(1)	O(1)
frontValue()	O(1)	O(1)

Applications of Stacks

◆ Balancing Parentheses

- ❑ Is this balanced?
 - $((()((()()())())))$
- ❑ Push when you see an open parenthesis.
- ❑ Pop when you see a closed one.
- ❑ If stack empties at end of string, it is balanced.
- ❑ Else, it is not.

◆ Postfix Expressions

- ❑ How to evaluate this:
 - $4\ 2\ * \ 5\ + \ 7\ + \ 1.5\ *$
- ❑ Push when you see a number
- ❑ Pop two when you see an operator, operate and push result

How to deal with (unary) negative signs such as $4 * (-5)$?

Applications of Stacks ... 2

◆ Infix to Postfix Conversion

□ How to convert this:

• **Infix:** $((4 * 2) + 5 + 7) * 1.5$ **to Postfix:** $4 2 * 5 + 7 + 1.5 *$

□ **Infix:** $4 * 2$

□ Idea: Push operators on a stack and print operands

□ **Infix:** $4 + 2 + 3$

□ Idea: Before pushing operators, pop previous operator

□ **Infix:** $4 + 2 * 5 + 3$

□ Idea: Before pushing operators, pop all operators of higher or equal precedence

□ **Infix:** $4 * (2 + 3) + 5$

□ Idea: Push open parenthesis on stack. When closed parenthesis is encountered, pop stack until open parenthesis is popped

Applications of Queues

- ◆ Servers

- Disk server, File Server, Print Server, ...