

### Evaluation

• Exams (2)	50%
• Homework Assignments	35%
• Semester Project	10%
• Class Participation	5%

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### Celebrity Problem

- A **Celebrity** is one that knows nobody and that everybody knows.

**Celebrity Problem:**  
 INPUT: n persons with a n×n information matrix.  
 OUTPUT: Find the “celebrity”, if one exists.  
 MODEL: Only allowable questions are:

- *Does person i know person j?*

- Naive Algorithm:  $O(n^2)$  Questions.
- Using Divide-and-Conquer:  $O(n \log_2 n)$  Questions.
- Improved solution?

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### Celebrity Problem (Cont'd)

- Naive Algorithm:  $O(n^2)$  Questions.
  - Ask everyone of everyone else for a total of  $n(n-1)$  questions
- Using Divide-and-Conquer:  $O(n \log_2 n)$  Questions.
  - Divide the people into two equal sets. Solve recursively and find two candidate celebrities from the two halves. Then verify which one (if any) is a celebrity by asking  $n-1$  questions to each of them and  $n-1$  questions to everyone else about them. This gives a recurrence for the total number of questions asked:  $T(n) = 2T(n/2) + 2n$
- Improved solution?
  - Hint: What information do you gain by asking one question?

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### Celebrity Problem (Cont'd)

- **Induction Hypothesis 2:** We know how to find  $n-2$  non-celebrities among a set of  $n-1$  people, i.e., we know how to find at most one person among a set of  $n-1$  people that could potentially be a celebrity.
- Resulting algorithm needs  $[3(n-1)-1]$  questions.

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### Solving Recurrence Relations

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Recurrence; Cond	Solution
$T(n) = T(n-1) + O(1)$	$T(n) = O(n)$
$T(n) = T(n-1) + O(n)$	$T(n) = O(n^2)$
$T(n) = T(n-c) + O(1)$	$T(n) = O(n)$
$T(n) = T(n-c) + O(n)$	$T(n) = O(n^2)$
$T(n) = 2T(n/2) + O(n)$	$T(n) = O(n \log n)$
$T(n) = aT(n/b) + O(n);$ $a = b$	$T(n) = O(n \log n)$
$T(n) = aT(n/b) + O(n);$ $a < b$	$T(n) = O(n)$
$T(n) = aT(n/b) + f(n);$ $f(n) = O(n^{\log_b a - \epsilon})$	$T(n) = O(n)$
$T(n) = aT(n/b) + f(n);$ $f(n) = O(n^{\log_b a})$	$T(n) = \Theta(n^{\log_b a} \log n)$
$T(n) = aT(n/b) + f(n);$ $f(n) = \Omega(n^{\log_b a})$ $af(n/b) \leq cf(n)$	$T(n) = \Omega(n^{\log_b a} \log n)$

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