

FALL 2008: **COT 5407** INTRO. TO ALGORITHMS  
[HOMEWORK 1; DUE SEP 23 AT START OF CLASS]

**General submission guidelines and policies:** ADD THE FOLLOWING SIGNED STATEMENT. Without this statement, your homework will not be graded.

I HAVE ADHERED TO THE COLLABORATION POLICY FOR THIS CLASS. IN OTHER WORDS, EVERYTHING WRITTEN DOWN IN THIS SUBMISSION IS MY OWN WORK. FOR PROBLEMS WHERE I RECEIVED ANY HELP, I HAVE CITED THE SOURCE, AND/OR NAMED THE COLLABORATOR.

Read the handout on **Homework guidelines and collaboration policy** from your course website before you start on this homework. This is very important.

## Problems

0. (**Regular**) Did you follow the instructions above?
1. (**Exercise**) Write down the time complexities of performing LINEARSEARCH and BINARYSEARCH in a sorted array of  $n$  elements.
2. (**Exercise**) Write down the precise **invariants** for each of the following algorithms: SELECTIONSORT, INSERTIONSORT, BUBBLESORT, SHAKERSORT, and MERGE (not MERGESORT).
3. (**Regular**) Given two monotonically increasing functions,  $f(n)$  and  $g(n)$ , prove or disprove:

$$\min(f(n), g(n)) = \Theta(f(n) + g(n)).$$

Redo the above problem with min replaced with max.

4. (**Regular**) (Exercise 4-1(c,d), p85) Use the *Master method* to solve the recurrences in (a), (b), and (c). You may assume that  $T(n)$  is constant for  $n < 2$ .
  - (a)  $T(n) = 15T(n/4) + 4n^2$ .
  - (b)  $T(n) = 27T(n/3) + 3n^3$ .
  - (c)  $T(n) = 18T(n/2) + 2n^4$ .
5. (**Regular**) Solve one of the problems in Problem 4 above using the *Substitution method*.
6. (**Regular**) Show that for any real constants  $a$  and  $b$ , where  $b > 0$ ,

$$(n - a)^b = \Theta(n^b).$$

Note that  $f(n) = \Theta(g(n))$  if and only if  $f(n) = O(g(n))$  and  $g(n) = O(f(n))$ .

7. (**Exercise**) (Exercise 3-2(e), p58) If  $f(n) = n^{\log_2 c}$  and  $g(n) = c^{\log_2 n}$ , indicate which of these relationships are true and prove your answers:  $f(n) = O(g(n))$ ,  $f(n) = \Omega(g(n))$ , and  $f(n) = \Theta(g(n))$ .
8. (**Extra Credit**) In our first class (Sep 2), we discussed and analyzed a simple algorithm for the SEARCH problem. We discussed two variants – one where  $X$ , the number to be searched, was bounded on both sides, and another where  $x$  was bounded below, but unbounded above. Binary search was the best strategy for the first version. The best strategy for the second version involved doing a doubling search followed by a binary search. This could be thought of as doing a LINEARSEARCH for  $m$ , the smallest exponent of 2 greater than  $x$ . What if we consider doing doubling search for  $m$ ? Can we push this even further? Analyze the best algorithm for this problem.