

COT 6405: Analysis of Algorithms

Giri NARASIMHAN

www.cs.fiu.edu/~giri/teach/6405F19.html

More Dynamic Operations

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Unsorted Arrays	$O(N)$	$O(1)$	$O(N)$	
Sorted Arrays	$O(\log N)$	$O(N)$	$O(N)$	
Unsorted Linked Lists	$O(N)$	$O(1)$	$O(N)$	
Sorted Linked Lists	$O(N)$	$O(N)$	$O(N)$	
Binary Search Trees	$O(H)$	$O(H)$	$O(H)$	$H = O(N)$
Balanced BSTs	$O(\log N)$	$O(\log N)$	$O(\log N)$	As $H = O(\log N)$

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Balanced BSTs	$O(\log N)$	$O(N)$	$O(N)$	
Augmented BBSTs	$O(\log N)$	$O(\log N)$	$O(\log N)$	

Room Scheduling Problem

- Given a set of requests to use a room
 - [0,6], [1,4], [2,13], [3,5], [3,8], [5,7], [5,9], [6,10], [8,11], [8,12], [12,14]
- Schedule largest number of above requests in the room
- Different approaches
 - Try by hand, exhaustive search, improve an initial solution, iterative methods, divide and conquer, greedy methods, etc.
- **Simple Greedy Selection**
 - Sort by start time and pick in “greedy” fashion
 - Does not work. WHY?
 - [0,6], [6,10] is the solution you will end up with.
- **Other greedy strategies**
 - Sort by length of interval
 - Does not work. WHY?

Room Scheduling – Improved Solution

- [0,6], [1,4], [2,13], [3,5], [3,8], [5,7], [5,9], [6,10], [8,11], [8,12], [12,14]
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]
-- Sorted by finish times
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]
- [1,4], [3,5], [0,6], [5,7], [3,8], [5,9], [6,10], [8,11], [8,12], [2,13], [12,14]

Greedy Algorithms

- Given a set of activities (s_i, f_i) , we want to schedule the maximum number of non-overlapping activities.
- **GREEDY-ACTIVITY-SELECTOR** (s, f)
 1. $n = \text{length}[s]$
 2. $S = \{a_1\}$
 3. $i = 1$
 4. for $m = 2$ to n do
 5. if s_m is not before f_i then
 6. $S = S \cup \{a_m\}$
 7. $i = m$
 8. return S

Why does it work?

➔ THEOREM

Let A be a set of activities and let a_1 be the activity with the earliest finish time. Then activity a_1 is in some maximum-sized subset of non-overlapping activities.

➔ PROOF

Let S' be a solution that does not contain a_1 . Let a'_1 be the activity with the earliest finish time in S' . Then replacing a'_1 by a_1 gives a solution S of the same size.

Why are we allowed to replace? Why is it of the same size?

Then apply induction! *How?*

Why does it work? Contd...

- **First choice was a good choice. Why?**
 - **Because it can be extended to an optimal soln.**
- **If our first choice was a good choice, then?**
 - **Then we can recursively apply correctness to the remainder**