

Query Optimization illustration

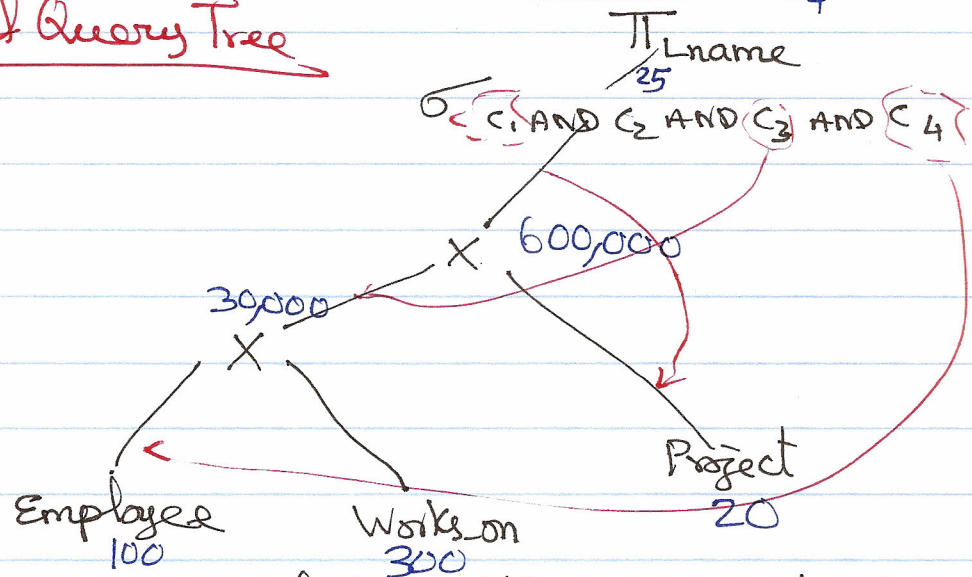
List the lastname of employees who work for 'AQUARIUS' Project and born after '1982-12-31'.

Given SQL Query

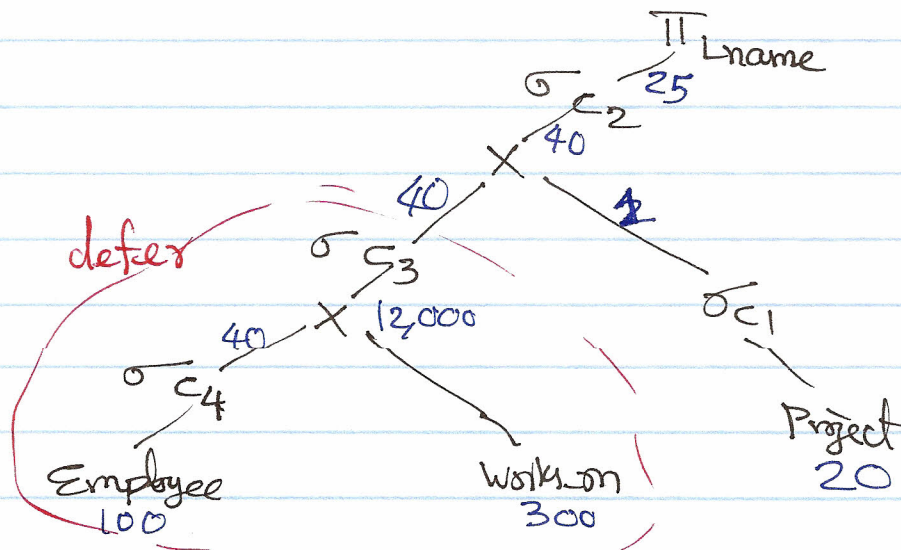
```

    ③ SELECT Lname
    ① FROM Employee, works_on, Project
    ② WHERE Pname = 'AQUARIUS' AND
           Pnumber = Pno AND
           ESSN = SSN AND
           Bdate > '1982-12-31'
    
```

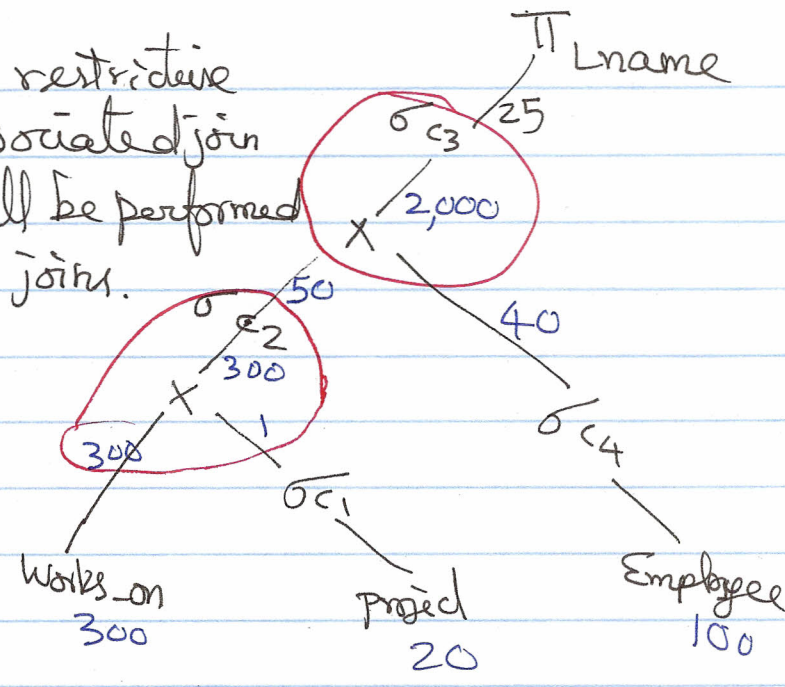
Initial Query Tree



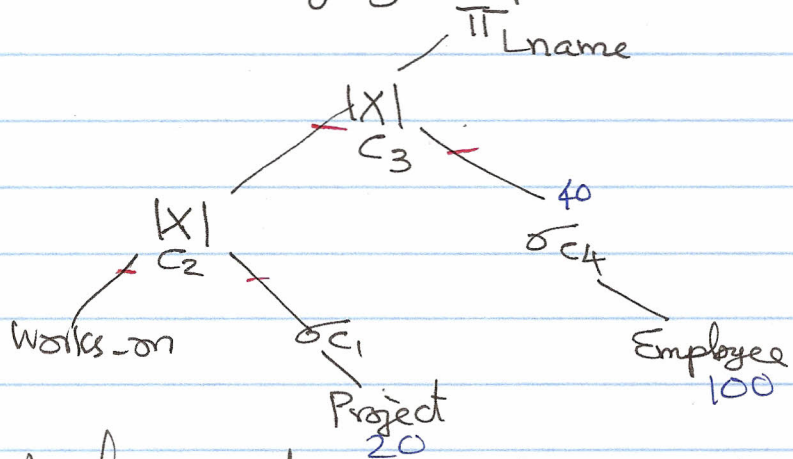
Step 1: Move select conditions down the tree.



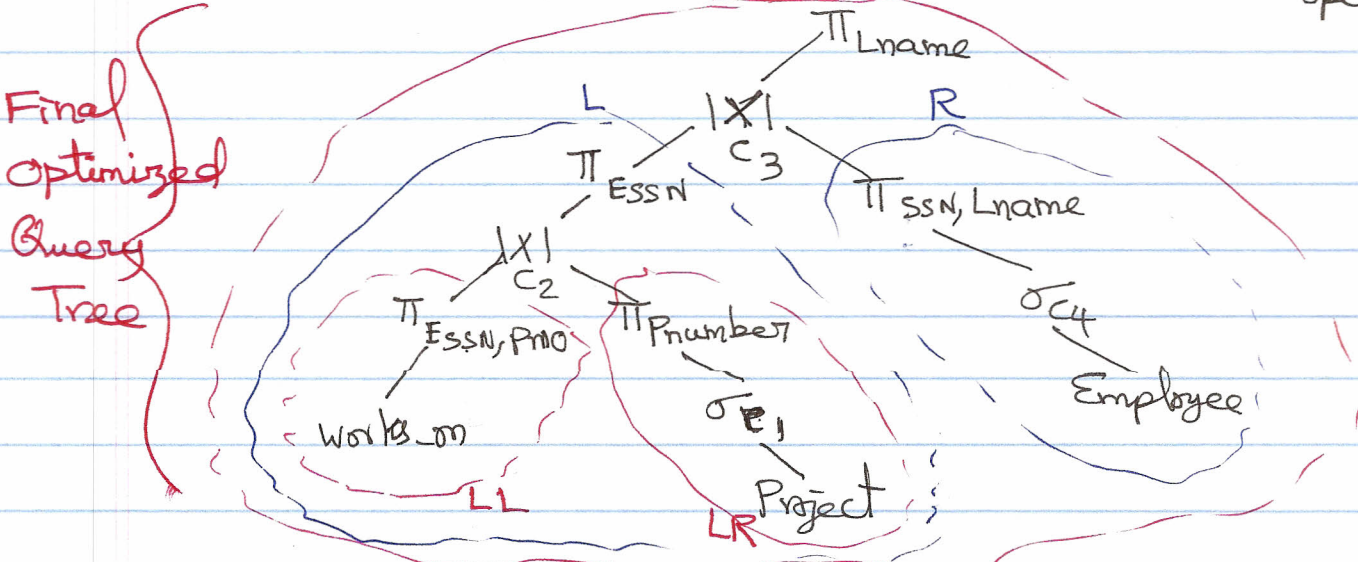
Step 2: Swap the most restrictive selection & its associated join so that this join will be performed earlier than other joins.



Step 3: Replace cross product & selection pair with sort-merge join operation



Step 4: Apply projection on each incoming branch of join operations.



```

SELECT Lname
FROM ( SELECT ESSN
      FROM ( LL )
      JOIN
      ( LR )
      ON C2 ) as L
JOIN
( R )
ON C3

```

Complete SQL Query for the optimized tree:

```

SELECT Lname
FROM ( SELECT ESSN
      FROM ( SELECT ESSN, PNO
            FROM works_on ) as LL
      JOIN
      ( SELECT Pnumber
        FROM Project
        WHERE Pname = 'AQUARIUS' ) as LR
      ON LL.PNO = LR.Pnumber ) as L
JOIN
( SELECT SSN, Lname
  FROM Employee
  WHERE Bdate > '1982-12-31' ) as R
ON L.ESSN = R.SSN;

```

Tuning Queries

- Partitioning^a table and store them on multiple physical disks to exploit concurrent disk I/O operation.
- Creating relevant indexes for retrieval queries.
- Apply deNormalization and store derived results to speed up future queries on the same set of tables.